

作业反馈3-4

TC练习**18.1.1**; 18.1.4; **18.2.3**; **18.2.4**; 18.3.1

18.1-1

Why don't we allow a minimum degree of $t = 1$?

18.2-3

Explain how to find the minimum key stored in a B-tree and how to find the predecessor of a given key stored in a B-tree.

- To find the minimum key in a B-tree
 - start at the root.
 - If the current node has children, then go to the leftmost child.
 - If the current node is a leaf, return the minimum key stored in the current node.

To find the predecessor of a given key, find the position of the given key in the tree and start at that node.

1. If the node is **not a leaf**, return the maximum key in the left subtree of the two subtrees that the key separates.
2. If the node **is a leaf** and the key is not the minimum key in the leaf, return the maximum key in the node that is less than the given key.
3. If the node is **a leaf and the key is the minimum key in the leaf**, go to the father of the node. If in the father node there is a key that is less than the given key, return it. Otherwise go to the father of this node and repeat this process.
4. If in the above process we **reach the root and still cannot find** the predecessor, then the given key is the minimum key in the tree.

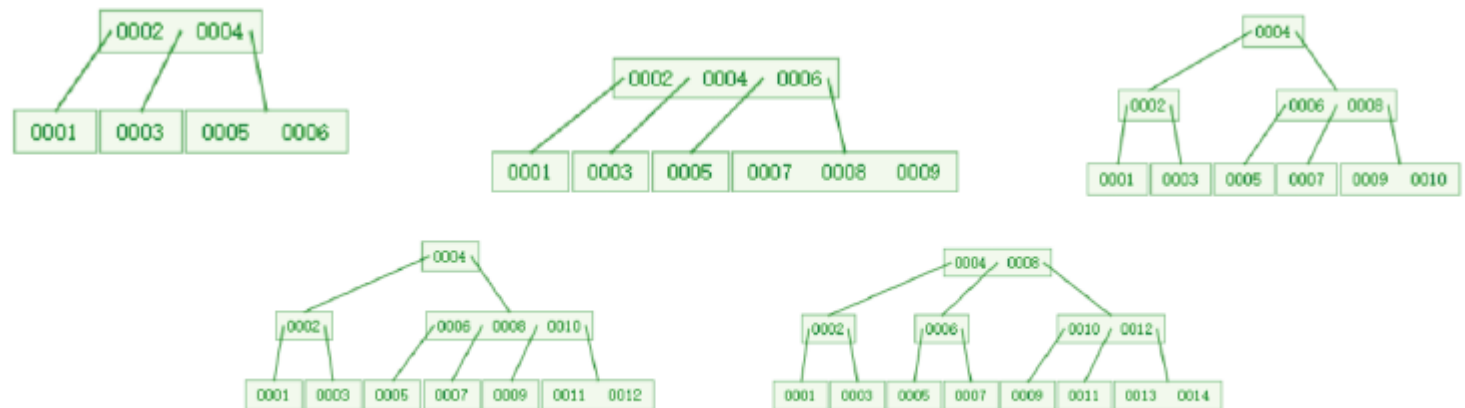
18.2-4 ★

Suppose that we insert the keys $\{1, 2, \dots, n\}$ into an empty B-tree with minimum degree 2. How many nodes does the final B-tree have?

- $\Theta(n)$?
 - 对，但是我不心甘！
 - 能否更加精确？

仅考虑 1-n 顺序插入的情况

由于 2-3-4 树的性质，在分裂时必定选择第 2 个 *key* 作为新的根节点并向上合并，这使得新节点的左边的部分只有一个 *key*，如图



并非书上介绍的single pass

直观上的说，如果插入是顺序的，那么只会改变整棵树右下角的节点，这使得整个树仅在最右支存在具有多个key的节点，并且根据分裂的性质，在最右支上，除了根节点可以有1 或2 或3个key，其他节点均有2 或3 个key

18.2-4 ★

Suppose that we insert the keys $\{1, 2, \dots, n\}$ into an empty B-tree with minimum degree 2. How many nodes does the final B-tree have?

Each node can have 1, 2, or 3 keys. Define the number of nodes that have at least 2 keys and 3 keys as a and b , respectively. Then the number of nodes is given by $n - a - b$.

Only the nodes on the rightmost path can have at least 2 keys, so $a = h + 1$ when the root has at least 2 keys and $a = h$ when the root has only 1 key, where h is the height of the tree.

- $f(h)$ as the number of keys when the nodes on the rightmost path except the root have two keys and other nodes have only one key

$$\begin{aligned} f(h) &= (2^h - 1) + 2(h + 1) - 1 + 2(2^h - 1 - h) \\ &= 3 * 2^h - 2 \end{aligned}$$

- $g(h)$ as the number of keys when all the nodes on the rightmost path have three keys and other nodes have one key.

$$\begin{aligned} g(h) &= 3(2^h - 1) + 3(h + 1) + 3(2^h - 1 - h) & g(h) &= 3(2^h - 1) + 3 + g(h) \quad \text{, when } h > 0. \\ &= 6 * 2^h - 3 & g(0) &= 3. \end{aligned}$$

18.2-4 ★

Suppose that we insert the keys $\{1, 2, \dots, n\}$ into an empty B-tree with minimum degree 2. How many nodes does the final B-tree have?

Since

$$f(h) \leq n \leq g(h),$$

we have

$$\log_2 \frac{n+3}{6} \leq h \leq \log_2 \frac{n+2}{3}.$$

For any integer n , the following inequality holds,

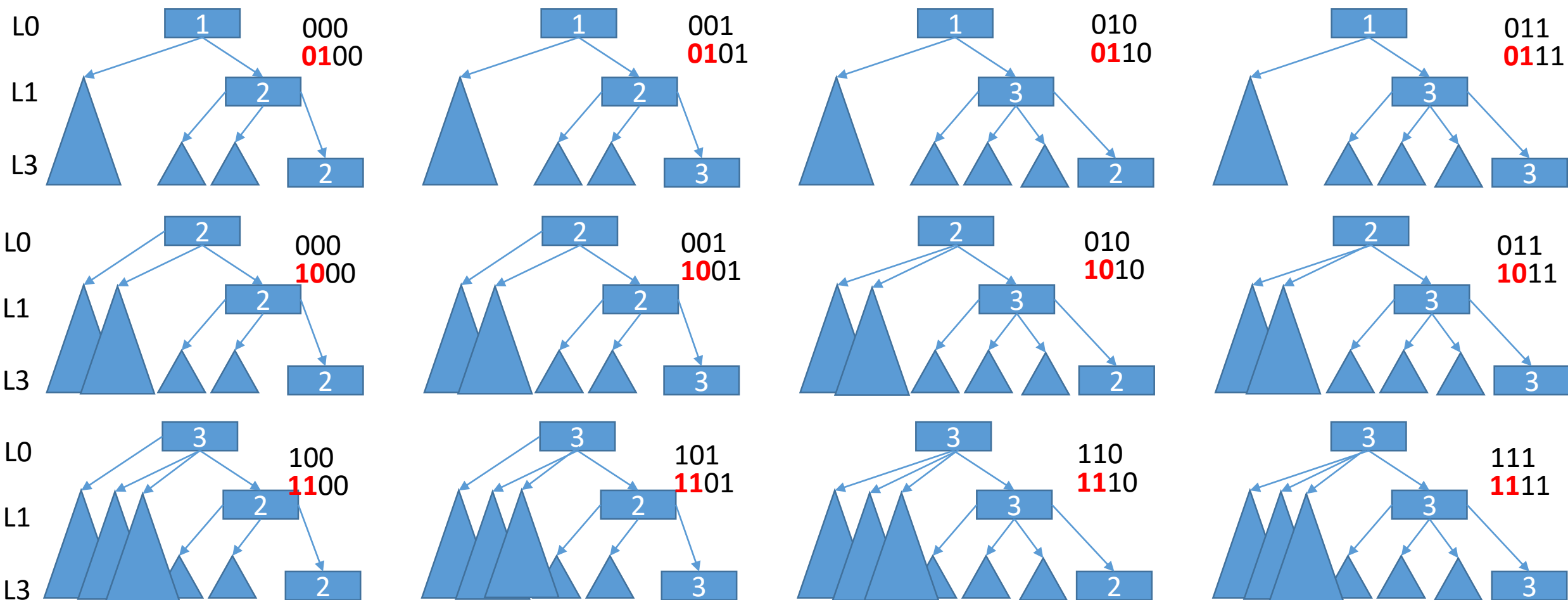
$$\log_2 \frac{n+2}{3} - \log_2 \frac{n+3}{6} < 1$$

Thus for any given integer n , there is only one integer h that satisfies the above inequality, which is $\lfloor \log_2 \frac{n+2}{3} \rfloor$.

18.2-4 ★

Suppose that we insert the keys $\{1, 2, \dots, n\}$ into an empty B-tree with minimum degree 2. How many nodes does the final B-tree have?

Notice that when $n = f(h)$, there are no nodes that have 3 keys. If one more key $n + 1$ is inserted, the rightmost leaf of the tree will have 3 keys. If $n + 2$ is inserted, then the rightmost leaf will split and the father of it will have 3 keys.



18.2-4 ★

Suppose that we insert the keys $\{1, 2, \dots, n\}$ into an empty B-tree with minimum degree 2. How many nodes does the final B-tree have?

- 对于按序插入 $1 \sim n$ ，高度为 h 的满足条件的B-tree
 - 我们可以将 $n - f(h)$ 表示为一个 $h+2$ 位的二进制串；
 - 则包含3个key的节点个数 b 可以表示为：
 - $b = (value_{1 \sim 2}(n - f(h)) \geq 3 ? 1 : 0) + bitcount_{3 \sim h+2}(n - f(h))$
 - 其中
 - $value_{1 \sim 2}(n - f(h))$: $n - f(h)$ 的二进制表示(强制 $h+2$ 位)前两位的值
 - $bitcount_{3 \sim h+2}(n - f(h))$: $n - f(h)$ 二进制表示(强制 $h+2$ 位)的 $3 \sim h+2$ 位中1的个数

$$\#Node(n) = n - a - b$$

18.3-1

Show the results of deleting C , P , and V , in order, from the tree of Figure 18.8(f).

